



THE PREMIER CONFERENCE & EXHIBITION ON COMPUTER GRAPHICS & INTERACTIVE TECHNIQUES

# Helix-Free Stripes for Knit Graph Design

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**Goal**: Given a 3D mesh, generate a **helix-free** knit graph over it

Must be helix-free!



<sup>1</sup> from Visual Knitting Machine Programming (2019) <sup>2</sup>from Knit Sketching: from Cut and Sew Patterns to Machine-Knit Garments (2021)

### Positioning our work

- Our goal (similar to Autoknit [Naryanan et al. 2018]):
  - Automatic generation of machine-knittable graphs from input triangle meshes
  - Precise control over helices
- 1-form-based framework produces more globally-informed stitch patterns

• Optimization allows incorporation of linear user-specified constraints









# Positioning our work

- We follow a stripes-based methodology
  - Evenly-spaced stripes ↔ evenly-spaced course rows and wale columns

- KnitKit [Nader et al. 2021] is only other work to consider this
  - They intersect stripes produced by [Knoppel et al. 2015], which often contain helices
  - Removal attempted via quad mesh operations [Bommes et al. 2011], but no guarantees
  - Our constraints can be used to guarantee helix removal





# $\rightarrow \sigma$ : The Stripe Texturing Function

 $\sigma_{ki}$ 



• 1 form  $\sigma: E \to \mathbb{R}$  a discretization of a vector field

 $\langle \sigma_{jk} \rangle$ 



- If  $(d_1\sigma)_m = 0$  vector field is integrable to a local linear function  $\circ$  Striping red if pixel value  $\in (0,\pi) \pmod{2\pi}$ , pink if value  $\in (\pi,2\pi) \pmod{2\pi}$
- If  $(d_1\sigma)_m = (2\pi)k, k \in \mathbb{Z}^{|F|}$ , vector field non-integrable

 $\sigma_{ij}$ 

- $\circ$  Can still get local function from triangle to  $\ S^1$
- Allows for global striping

•

### $\rightarrow$ Optimizing for $\sigma$ directly

- We optimize directly for  $\,\,\sigma\,\,$ 
  - [Noma et al. 2022] optimize in the space of 1-forms but do not explore knitting applications

- The gradient of a harmonic interpolation  $h: V \to \mathbb{R} \in [0, 1]$  guides the 1-form optimization
  - Linear constraints achieve all the desiderata

$$\underbrace{||W(\sigma_c - \omega_c)||^2}_{\text{Quadratic minimization objective for course rows}}$$

$$(\omega_c)_{ij} = \underbrace{\frac{1}{2} \left( \frac{(\nabla h)_l}{||(\nabla h)_l||} + \frac{(\nabla h)_r}{||(\nabla h)_r||} \right) \cdot \mathbf{e}_{ij}}_{\text{integral of normalized } \nabla h \text{ along } e_{ii}}$$











Stripe irregularity/short row placement



Direct use of [Knoppel et al. 2015]

Tightness to boundaries through stripe alignment

$$\sigma_c|_{\partial M} = 0$$



 $(d_1\sigma)_m = (2\pi)k, k \in \mathbb{Z}^{|F|}$ 







This level of granular control not achievable with linear constraints in [Knoppel et al. 2015]



### Optimization problems and strategies









Thanks to Yuxuan Mei in Adriana Schulz's group at UW for help with fabricating!







- Presented a **stripes-based framework** for global stitch structure specification, operating in the space of **discrete 1-forms**
- Results in quadratic mixed-integer optimization with **linear** constraints for:
  - Removal of helices
  - Placement of stitch irregularities
  - Alignment of course rows and wale columns to feature curves
- Constraints are directly incorporated into a **global structure optimization** 
  - Most previous works incorporate user constraints via post-processing



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# **Thank you for listening!**

(Slides that follow are for clarifying question)

# Why do knit graphs need to be helix-free?



- Knitting schedulers [Naryanan et al. 2018] trace the graph in a natural spiralling/helical pattern
  - Helices in the generated graph would result in an incorrect amount (i.e., too much) helicing





<sup>1</sup>from Automatic Machine Knitting of 3D Meshes(2018)







# Level Set (path integral constraints)

- Can specify path integral of  $\,\sigma\,$  along a specified polyline  $\,\gamma\,$ 
  - Constrains same stripe level set to pass through polyline endpoints
  - Non-trivial constraints applied to the wale stripes to ensure stripes "close up" on the boundary

$$\int_{\gamma} \sigma = \sum_{i=0}^{n-1} -s_i^{i+1} \sigma_{e_i} r_i^{i+1} + s_{i+1}^i \sigma_{e_{i+1}} r_{i+1}^i = C, C \in \mathbb{R}$$

• When C = 0 and  $\gamma$  is a closed loop of mesh edges, ensures that a stripe level set does not cross the loop without returning back to its side of origin i.e., helix removal











### Constraints in the wale direction



$$(\omega_w)_{ij} = \frac{1}{2} \left( \frac{(\nabla h)_l^T}{||(\nabla h)_l^T||} + \frac{(\nabla h)_r^T}{||(\nabla h)_r^T||} \right) \cdot \mathbf{e}_{ij}$$

$$\int_{(\partial M)^i} \sigma_w = (2\pi) k^i$$
 Specified on N-1 boundaries

- Need to only consider irregularity placement and possible stripe alignment constraints
  - No helices in the wale direction









Model	Strategy	#V	#F	# Int.	MIP Solve
				Vars.	Time, s
Curved	S2	54	96	97	< 5
Cylinder (Fig. 7)					
Hemisphere	S2	224	384	385	40
(Fig. 9(e))					
Sock (Fig. 1)	S2	279	538	98	21
Sock	S1	279	538	2	< 1
(Fig. 9(d))					
Cylinder	S1	270	580	3	< 1
(Fig. 6, middle)					
Cactus	S2	391	736	738	35
(Fig. 9(c))					
Bunny	S2	2669	5228	5230	287
(Fig. 4, right)					

#### Table 1: Run time statistics.

